### **SECURITY, PERFORMANCE: PICK ONE ?**

Pierre Chifflier @pollux7

### WHAT

- Defensive/Secure programming
- Needed for many programs
  - Web servers, IDS
- Often rejected for perf reasons
- We want performance **and** security
- And something we can maintain over time

## WHO

- Debian Developer, Suricata contributor
- Head of the Detection Research lab at ANSSI
- Security, compilers and languages
- Write (rust) parsers for everything







### DISCLAIMER

- This talk is my story, not a guide
  - Could be "story of a fight vs the compiler"
- Still, there are general rules/hints

## **CHOOSING A LANGUAGE**

- Sometimes C is not the answer
  - You will *always* fail somewhere
  - Believe me, l've tried
- Parts in assembly
  - Could be fast, but a nightmare to maintain
- OCaml, Go, ...
  - The perf-killer garbage collector

### **EXAMPLE: PARSING KERBEROS**

KDC-REQ-BODY := SEQUENCE {		
kdc-options	[0] K	DCOptions,
cname	[1] P	rincipalName OPTIONAL
	-	- Used only in AS-REQ,
realm	[2] R	ealm
	-	- Server's realm
	-	- Also client's in AS-REQ,
sname	[3] P	rincipalName OPTIONAL,
from	[4] K	erberosTime OPTIONAL,
till	[5] K	erberosTime,
rtime	[6] K	erberosTime OPTIONAL,
nonce	[7] U	Int32,
etype	[8] S	EQUENCE OF Int32 EncryptionType
	-	- in preference order,
addresses	[9] H	ostAddresses OPTIONAL,
enc-authorization-data	[10]	EncryptedData OPTIONAL
	-	- AuthorizationData,
additional-tickets	[11]	SEQUENCE OF Ticket OPTIONAL
		NOTE: not empty

### **ASN.1 DER ENCODING**

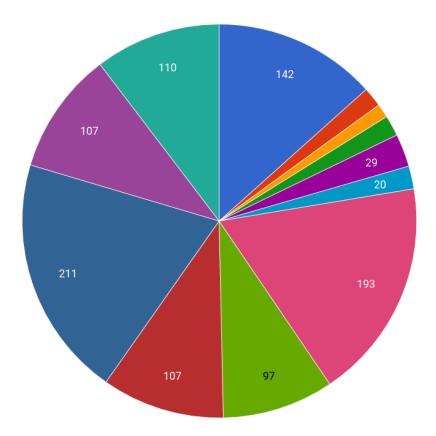
#### The 9 layers of DER Hell

- Lots of TLVs
- Highly recursive
- Infinite size integers
- Variable lengths





### WRITING PARSERS IS HARD



- heap buffer overflows
- global buffer overflows
- stack buffer overflows
- use after frees
- uninitialized memory
- stack overflows
- timeouts
- ooms
- leaks
- ubsan
- unknown crashes
- other (e.g. assertions)

(source: Google OSS-Fuzz)

# RUST



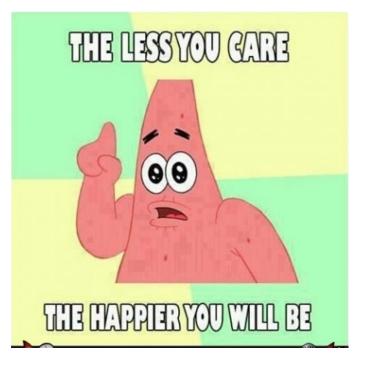
- Compatible with C
- Type safety
- Memory safety: non uninitialized values, etc.
- Thread safety: forces you to protect (lock) concurrent access
- Note: integer overflow/underflow still possible

### **RUST SPEED**

- Fast, but not enough
- Can we do better
  - Keeping safety
  - Keeping some readability

### MANDATORY WARNING

- First rule: don't optimize
- Don't bother one-time optimizations



# METHODOLOGY

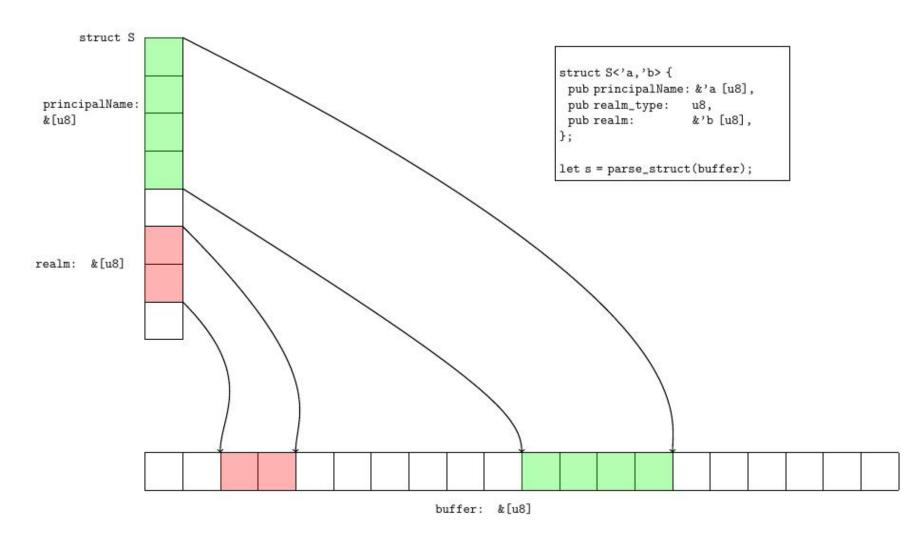
- First: identifying slow points
- Use available tools
  - cargo bench
  - perf, valgrind
  - flamegraph
- One eye on the source code
- Also look at the produced binary code



# **ACTION 1: SOURCE CODE**

- Algorithms first
- Zero-copy
  - Made possible thanks to slices
- Non-locking code
  - Borrow-checker and non-mutability help *a lot*

### **SLICES**



### **RESULT 1**

- messages of 305 bytes
- 1856 ns / message -> 156 MB/s (per thread)
- Fast, but we want more

### **ADDING INSTRUMENTATION: STEP 1**

#### • Add to Cargo.toml:

[profile.release]
debug = true

[profile.bench]

debug = true

### **STEP 2: ADD BENCHMARKS**

#### • Add a benchmark (benches/b\_krb5\_parser.rs):

```
static KRB5_TICKET: &'static [u8] = include_bytes!("../assets/krb5-ticket.bin");
#[bench]
fn bench_parse_ticket(b: &mut Bencher) {
    b.iter(|| {
        let res = parse_krb5_ticket(KRB5_TICKET);
        // use result !
        match res {
            Ok((rem,tkt)) => {
                assert!(rem.is_empty());
                assert_eq!(tkt.tkt_vno, 5);
                },
                __=> assert!(false),
                }
            });
}
```

# STEP 3: COLLECT INSTRUMENTATION RESULTS

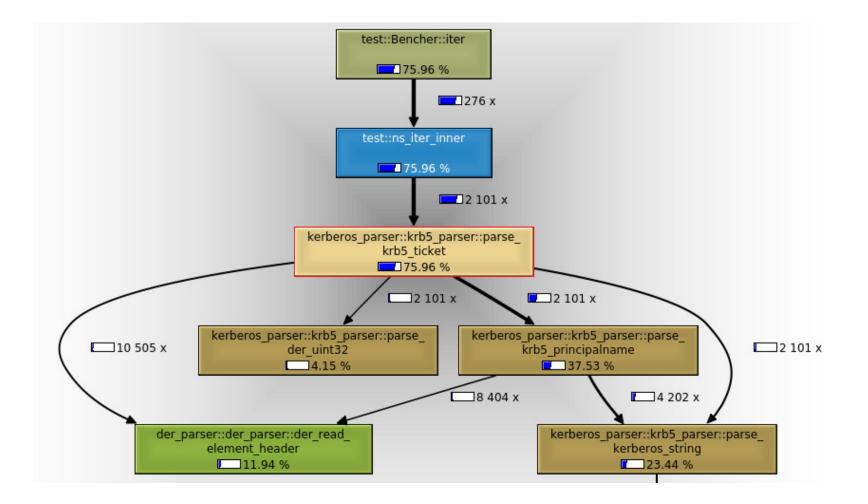
• Build the benchmark executable

cargo bench --no-run

#### • Use valgrind on it:

valgrind --tool=callgrind \
 --dump-instr=yes --collect-jumps=yes --simulate-cache=yes \
 ./target/release/b\_krb5\_parser-e84a853b88e37bef --bench bench\_parse\_ticket

#### **PERFORMANCE GRAPHS**



### WHY IS IT SLOW ?

- Parts of the code are slow
  - too many tests
  - useless data copy
- Some structures do not fit in cache



## ACTION 2: LOOK AT PRODUCED CODE

- Goal: try to identify and use efficient patterns
- Can bring huge speed improvements
- Time consuming
- Hard to find stable optimization patterns

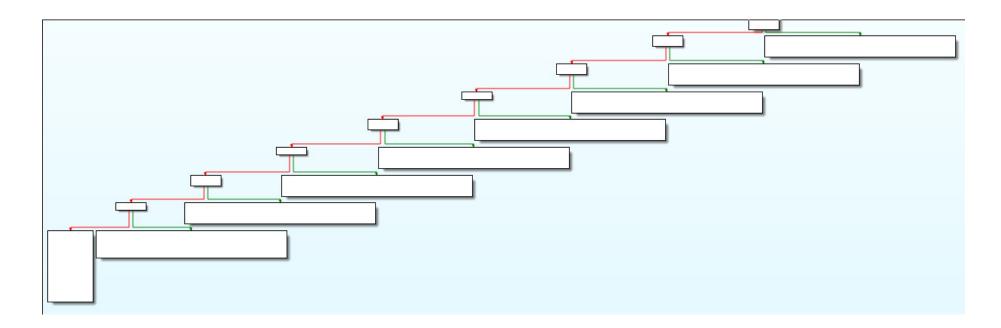
### **COMPILER IS YOUR FRIEND .. OR NOT**



#### **EXAMPLE: READING AN U64**

```
let u = (i[0] as u64) << 7 |
  (i[1] as u64) << 6 |
  (i[2] as u64) << 6 |
  (i[3] as u64) << 5 |
  (i[3] as u64) << 4 |
  (i[4] as u64) << 3 |
  (i[5] as u64) << 2 |
  (i[6] as u64) << 1 |
  (i[7] as u64);</pre>
```

#### **EXAMPLE: READING AN U64**



8 length tests

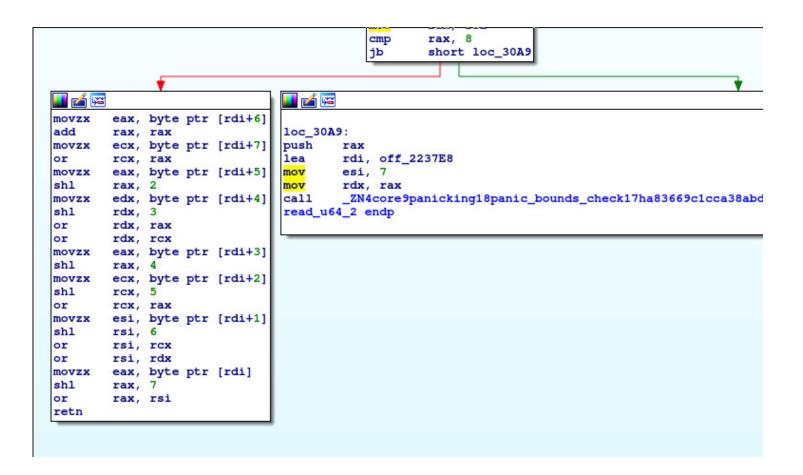
### **TEST #1: REORDERING**



(i[7]	as	u64)			
(i[6]	as	u64)	<<	1	
(i[5]	as	u64)	<<	2	
(i[4]	as	u64)	<<	3	
(i[3]	as	u64)	<<	4	
(i[2]	as	u64)	<<	5	
(i[1]	as	u64)	<<	6	
(i[O]	as	u64)	<<	7;	

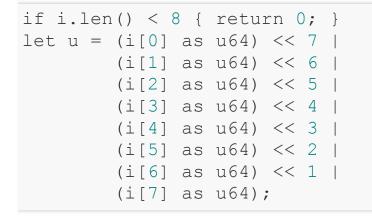
Read last bytes first

### **TEST #1: REORDERING**



Better, but we still have a panic statement

### **TEST #2: ASSERTING/TESTING SIZE**



- Compiler uses the info from the test
- No need to reorder

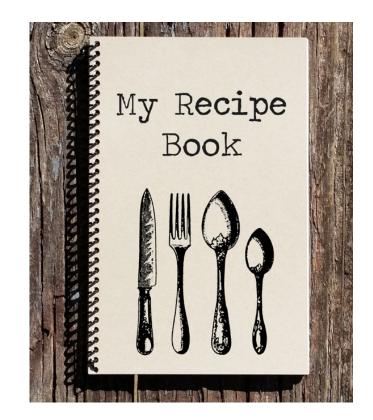
### **TEST #2: ASSERTING/TESTING SIZE**

;unwind {						
-	rsi, 7 short loc_30F9					
Ju .	shore rec_sers					
-						
💵 🖆 🖼						
xor eax, eax	1 2070.					
retn	loc_30F9:					
	movzx eax, byte ptr [rdi+6] add rax, rax					
	add rax, rax movzx ecx, byte ptr [rdi+7]					
	or rcx, rax					
	movzx eax, byte ptr [rdi+5]					
	shl rax, 2					
	movzx edx, byte ptr [rdi+4]					
	shl rdx, 3					
	or rdx, rax					
	or rdx, rcx					
	movzx eax, byte ptr [rdi+3]					
	shl rax, 4					
	movzx ecx, byte ptr [rdi+2]					
	shl rcx, 5					
	or rcx, rax					
	movzx esi, byte ptr [rdi+1]					
	shl rsi, 6					
	or rsi, rcx					
	or rsi, rdx					
	movzx eax, byte ptr [rdi]					
	shl rax, 7					
	or rax, rsi					
	retn					
	; } // starts at 30F0					
	read_u64_3 endp					
	L					

# **LESSONS FROM THAT EXAMPLE**

- Compiler is smart
- But not enough to infer all information
- Sometimes *adding* code makes result faster
  - Some tests/asserts have to be *explicit*
- We can get efficient code without using unsafe or assembly

### **SOME OTHER TIPS**



### PACKED ENUMS

Representing a packed enum:

#[repr(u8)]
pub enum Foo {
 Value1 = 1,
 Value2 = 2,
 ...

However:

- match is slow
- conversions to/from u8 are implemented as either
  - function calls (slow)
  - memory casts (unsafe)

### THE NEWTYPE PATTERN

pub struct Foo(pub u8);

- Type-safe, cost-free abstraction
- Free conversions
  - except if you forget the pub keyword!
- Compile time increases
- Values have to be declared as associated constants

impl Foo {
 pub const Value1 : Foo = Foo(1);

## ALLOCATIONS

- Allocations are slow
- Prefer the stack
  - Avoid Box and Vec
  - You can use variable-length data-types on stack
  - Drawback: calls to memcpy

### STRUCTURES

- Keep as much as possible in cache
  - Use small structs
  - Make sure they fit in cache
- Check using valgrind

## CODE

- Keep as much as possible in cache
- Keep as much as possible in registers
- Use reentrant, pure functions (no side-effects)
- Avoid locks and global structures
  - locks are slow!

### CODE

- Write *linear* code
  - Avoid instructions cache misses
- *nom* helps a lot (macros)
  - Possible problem: cyclomatic complexity

warning: the function has a cyclomatic complexity of 231
 --> src/krb5\_parser.rs:303:1

### **AUTOMATIC VECTORIZATION**

```
let len = min(min(a.len(), b.len()), c.len());
for i in 0..len {
    c[i] = a[i] + b[i];
}
```

#### Code is not vectorized:

```
rsi, r10
cmp
jae
        .LBB0 7
        rsi, rcx
cmp
jae
       .LBBO 8
        rsi, r9
cmp
jae
        .LBBO 9
        eax, dword ptr [rdi + 4*rsi]
mov
        eax, dword ptr [rdx + 4*rsi]
add
        dword ptr [r8 + 4*rsi], eax
mov
lea
        rax, [rsi + 1]
        rsi, rax
mov
. . .
```

### **AUTOMATIC VECTORIZATION**

```
let len = min(min(a.len(), b.len()), c.len());
let (a,b,c) = (&a[..len], &b[..len], &mut c[..len]);
for i in 0..len {
    c[i] = a[i] + b[i];
}
```

#### Code is vectorized:

```
movdqu xmm0, xmmword ptr [r10 + 4*rsi]
movdqu xmm1, xmmword ptr [r10 + 4*rsi + 16]
movdqu xmm2, xmmword ptr [rdx + 4*rsi]
paddd xmm2, xmm0
movdqu xmm0, xmmword ptr [rdx + 4*rsi + 16]
paddd xmm0, xmm1
...
```

#### However:

• Make sure instructions apply to blocks of data

## **MISC PATTERNS**

- Avoid Box<Trait>, prefer &mut Trait
  - former is 2 pointers (and extra checks)
- Use iterators
  - they can spare some more bounds checks

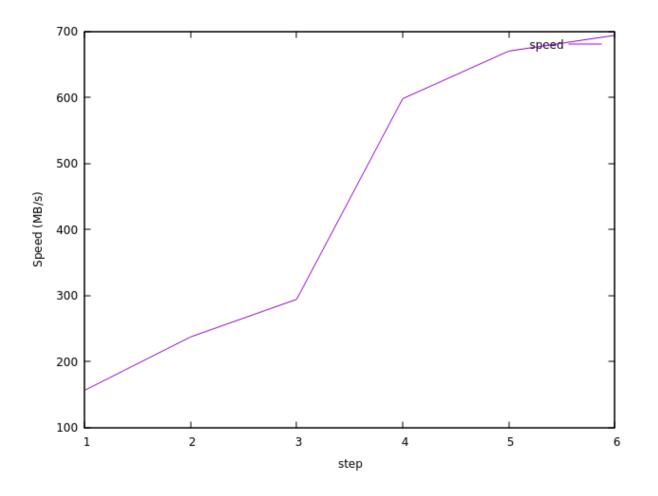
### LINKER

- Use LTO (Link-time optimization)
  - Bigger executable, generally faster
  - Better than using #[inline(always)]
- Test PGO (Profile-guided optimization)
  - Variable results
  - Sometimes really good

### **CHECKING WE STILL HAVE SECURITY**

- Compiler is not (always) your friend
  - Undefined behaviors
    - Integer underflows/overflows
    - Removed in release mode, but can be added
  - Removed tests
  - Removed calls (e.g. memset)
  - panic / assert inserted or remaining
- Use Compiler Explorer (https://rust.godbolt.org/)
- Use cargo fuzz

### **RESULTS (KERBEROS)**



• Comparison: my C implem. is ~ 500 MB/s

### CODE

- Kerberos, SNMP, IKEv2, TLS, Radius, etc.
- Rusticata project (and all parsers): https://github.com/rusticata
- Rust code push on crates.io
- Most parsers merged in Suricata git repo (4.1?)

### LESSONS

- Guiding the compiler is efficient
- And easier to maintain than writing assembly
- Always check that the tests are present after optimizing
- Are the optimizations stable ?
  - Not guaranteed, but in practise yes
  - Sometimes look like voodoo

